

Meshes

- polygonal soup
 - polygons specified one-by-one with no explicit information on shared vertices
- polygonal nonmanifold
 - connectivity information is provided (which vertices are shared)
no restrictions on connections between polygons
- polygonal manifold
 - no edge is shared by more than two polygons; the faces adjacent to a vertex form a single ring (incomplete ring for boundary vertices)
- triangle manifold
 - in addition, all faces are triangles

Mesh elements

faces, vertices, edges

Each mesh element can have information associated with it; typical mesh operations involve visiting (traversing) all vertices, faces, or edges

Mesh descriptions

- OBJ format
each line defines an element (vertex or face); first character defines the type

Vertex:

v x, y z

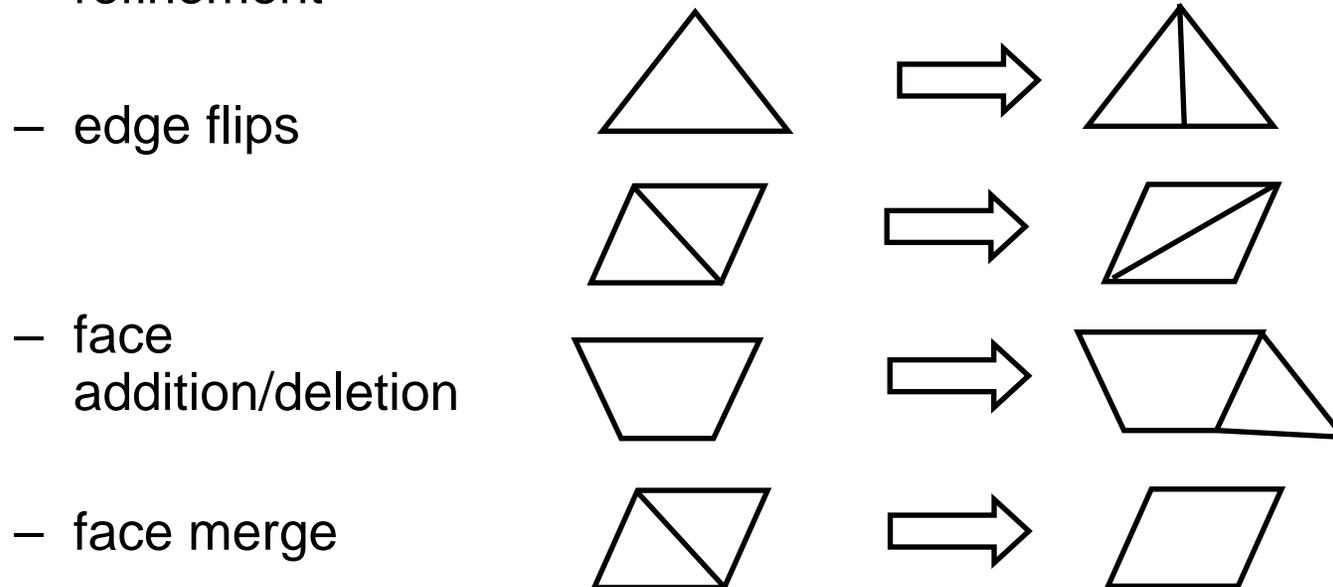
Face with n vertices:

f i1 i2 i3 ... in

where $i1.. in$, are vertex indices; the indices are obtained by numbering all vertices sequentially as they appear in a file

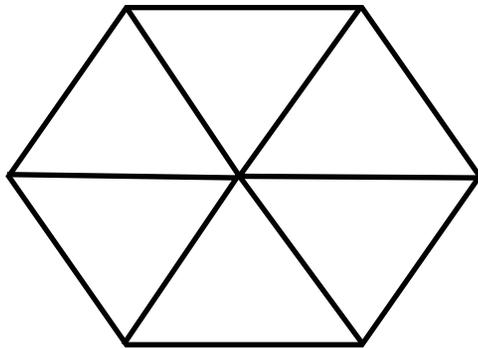
Mesh operations

- Types of mesh operations
 - traversals go over all elements of certain type
 - collect adjacent elements (e.g. all neighbors of a vertex)
 - refinement



Traversal operations

- Iterate over all vertices, faces, edge
 - visit each only once
 - iterate over all elements (faces, vertices, edges) adjacent to an element



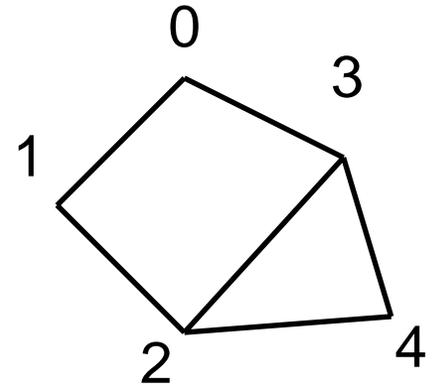
A simple mesh representation

One-to-one correspondence with OBJ

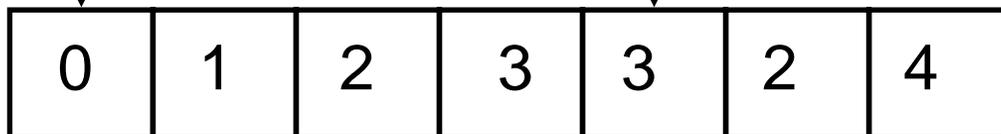
array of vertices

2 arrays for faces

each face is a list of vertex indices
enumerated clockwise



starting indices of face
vertex lists



vertex indices of
all faces

Traversal operations

Complexity of traversal operations w/o additional data structures as function of the number of vertices, assuming constant vertex/face ratio

iterate over collect adjacent	V	E	F
V	quadratic	quadratic	linear
E	quadratic	quadratic	linear
F	quadratic	quadratic	linear

Traversal operations

Most operations such as collecting all adjacent faces for a vertex are slow, because the connectivity information is not explicit: one needs to search the whole list of faces to find faces with a given vertex; if neighbors are encoded explicitly this can be done in const. time

Face-based mesh representation

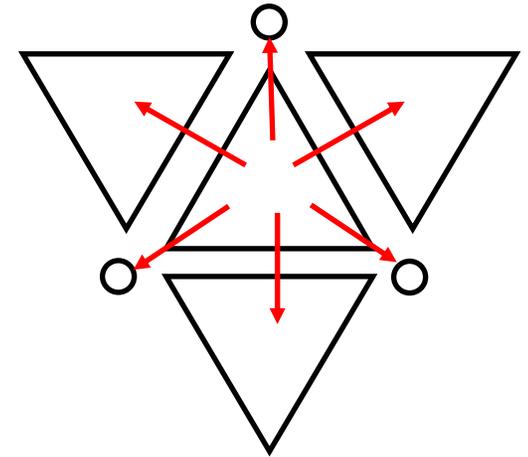
Useful primarily for triangle or quad. meshes

Triangle meshes:

```
struct Face {  
    Face*  face[3]; // pointers  
                                //to neighbors  
    Vertex* vertex[3];  
}
```

```
struct Vertex {  
    Face* face; // pointer to a triangle  
                //adjacent to the vertex  
}
```

(not really necessary, can refer to vertices using a handle (Face ptr, vertex index))



Traversing faces sharing a vertex

Assuming a mesh without boundary:

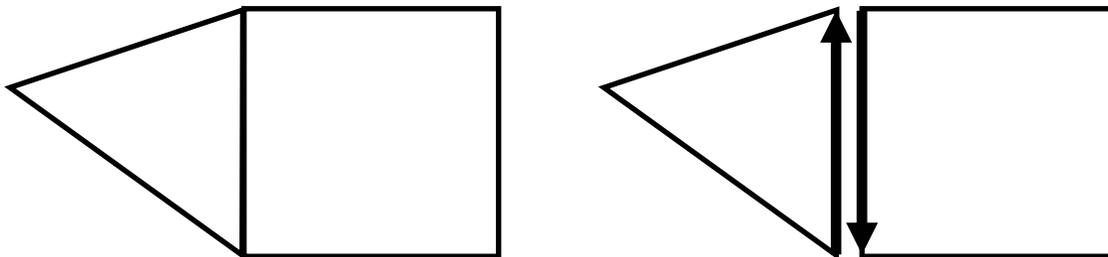
```
fstart = v->face;
f = fstart;
do {
    ... // perform operations with *f
    // assume that vertex i is across edge i
    if (f->vertex[0]== v)
        f = f->face[1]; // crossing edge #1  vert. 0 - vert. 2
    else if (f->v[1] == v)
        f = f->face[2]; // crossing edge #2  vert. 1 - vert. 0
    else
        f = f->face[0]; // crossing edge #0  vert. 2 - vert. 1
} while( f != fstart);
```

Similar for edges and vertices.

All such operations can be done in const. time per vertex/face/edge.

Half-edge data structure

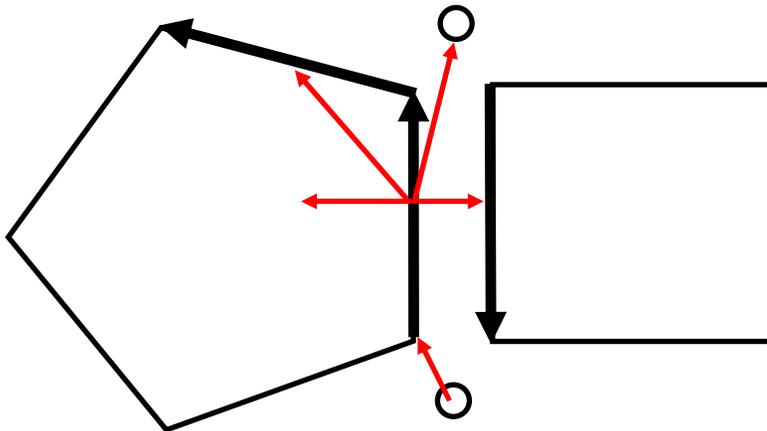
- General manifold polygonal meshes
 - Polygons have variable number of vertices variable size;
 - data structures based on faces are inconvenient and inefficient.
- Solution: use edge-based structures (winged edge, half-edge).
 - Half-edge is currently most common
 - Each edge = 2 half edges; can be interpreted either as
as
directed edge or face-edge pair



Half-edge data structure

```
struct HalfEdge {
    Vertex* vertex; // the head vertex the
                    // half edge is pointing to
    Face* face;    // if data stored in faces
    HalfEdge* next; // next halfedge in the face
                    // on the left
    HalfEdge* sym; // the other half edge for
                    // the same edge
}

struct Vertex {
    HalfEdge* halfedge; // one of the half edges
                        // starting at the vertex
}
```



Traversal operations

Vertices adjacent to a vertex v , mesh without boundary

```
he = v->halfedge;  
do {  
    he = he->sym->next;  
    ... // perform operations with  
        // he->vertex  
} while (he != v->halfedge)
```

No “if” statements.

Constructing a mesh data structure

Construct face-based structure from a list of triangles and vertices

Assume that vertices are listed counterclockwise for each triangle and v_i indices of vertices in the face; $\text{other}(i_1, i_2)$ for $i_1, i_2 = 0..2, i_1 \neq i_2$ is the third vertex of the triangle $i_3 \neq i_1, i_2$

Edgemap is a map (associative array) from pairs of vertices (directed edges) to faces; in addition to the face, we also record the number of the edge in the face (See C++ STL `map` details of use)

This is pseudocode (not using C syntax to emphasize this)

```
for each face
  create face structure f1, initialize neighbors to 0
  for each triangle vertex i=0..2
    edgemap(v_i, v_{(i+1)%3}) := (f1, other(i, (i+1)%3) )
  endfor
endfor
for each entry (i,j) of the map edgemap
  edgemap(i,j)
  (f2,e2) := edgemap(j,i);
  if f2 != 0 then
    f1->f[e1] := f2
    f2->f[e2] := f1
  endif
endfor
```

Building a half-edge data structure

- Similar to building face-based triangular mesh
- Input: a list of vertices, a list of faces, each face is a list of vertex indices enumerated CCW
- 1. Create arrays of vertices, faces and halfedges, one half-edge for every seq. pair of vertices of every face; initialize all pointers to zero.
- 2. For each face f , with n vertices
 - assign $f.\text{halfedge}$ to its first half-edge;
 - for each vertex v of a face, assign $v \rightarrow \text{halfedge}$ to the halfedge starting at it if nothing is assigned to it yet;
 - for each half-fedge he of a face, assign
 - $he.\text{face} = f$, $he \rightarrow \text{next}$ = next half-edge in the face,
 - $he \rightarrow \text{vertex}$ = next vertex in the face;
 - record half-edge pointer he in the edge map:
 - $\text{edgemap}(v[i], v[i+1]) = he$
- 3. Go over all entries of the edge map, assign for half-edges $\text{edgemap}(i, j)$ $\text{edgemap}(j, i)$ links to each other, if both exist

Dealing with boundaries

- To minimize implementation effort it is useful to create two halfedges for boundary edges, one of which has zero face pointer;
- A boundary vertex v should always have $v.\mathbf{halfedge}$ pointing to a boundary halfedge.
- Then it is easy e.g. to find two boundary neighbors of a vertex.